

MK68590 (P,N)

LOCAL AREA NETWORK

CONTROLLER FOR ETHERNET

COMMUNICATIONS PRODUCTS

FEATURES

- □ 100% compatible with Ethernet and IEEE 802.3 specifications
- Data packets moved by block transfers over a processor bus (on-board DMA controller 24-bit linear address space)
- □ Buffer management
- □ Packet framing
- ☐ Preamble and Cyclic Redundancy Check (CRC) insertion
- ☐ Preamble stripping and CRC verification
- General 16-bit microprocessor bus interface compatible with popular processors (68000, 8086, Z8000, LSI-11)
- □ Cable fault detection
- ☐ Multicast logical address filtration
- □ Collision handling and retry
- ☐ Scaled N-channel MOS VLSI technology
- ☐ 48-pin DIP
- ☐ Single 5-volt power supply
- ☐ Single phase TTL level clock
- ☐ All inputs and outputs TTL compatible
- ☐ Completely compatible with companion Serial Interface Adapter (SIA) chip MK68591/2.

DESCRIPTION

The MK68590-LANCE™ (Local Area Network Controller for Ethernet) is a 48-pin VLSI device that simplifies the interfacing of a microcomputer or a minicomputer to an Ethernet Local Area Network. This chip operates in a local environment that includes a closely coupled memory and microprocessor. The LANCE uses scaled N-channel MOS technology and is compatible with several microprocessors.

LANCE is a trademark of Thomson Components - Mostek Corporation.

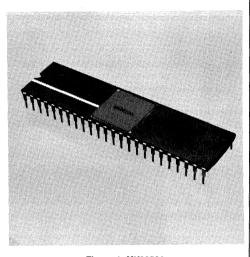


Figure 1. MK68590

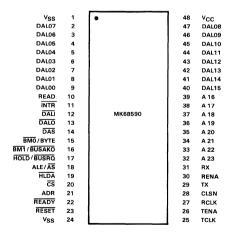


Figure 2. LANCE Pin Assignment

PIN DESCRIPTION

DAL00-DAL15

(Data/Address Bus)

Input/Output Three State. The time multiplexed Address/Data bus. These lines will be driven as a bus master and as a bus slave.

READ

Input/Output Three State. Indicates the type of operation to be performed in the current bus cycle. When it is a bus master, LANCE drives this signal.

LANCE as bus slave:

High - The chip places data on the DAL lines.

Low - The chip takes data off the DAL lines.

LANCE as bus master:

High - The chip takes data off the DAL lines. Low - The chip places data on the DAL lines.

INTR

(Interrupt)

Output Open Drain. When enabled, an attention signal that indicates the occurrence of one or more of the following events: a message reception or transmission has completed or an error has occurred during the transaction; the initialization procedure has completed; or a memory error has been encountered. Setting INEA in CSR0 (bit 06) enables INTR.

DALI

(Data/Address Line In)

Output Three State. An external bus transceiver control line. When LANCE is a bus master and reads from the DAL lines, DALI is asserted during the data portion of the transfer.

DALO

(Data/Address Line Out)

Output Three State. An external bus transceiver control line. When LANCE is a bus master and drives the DAL lines, DALO is asserted during the address portion of a read transfer or for the duration of a write transfer.

DAS

(Data/Strobe)

Input/Output Three State. Defines the data portion of the bus transaction. DAS is driven only as a bus master.

BM0, BM1 or BYTE, BUSAKO

(Byte Mask)

Output Three State. Pins 15 and 16 are programmable through bit (00) of CSR3 (known as BCON). Asserting RESET clears CSR3.

fBCON = 0

PIN 16 = BM1 (Output Three State)

PIN 15 = BM0 (Output Three State)

BM0, BM1 Byte Mask. Indicates the byte(s) of a bus transaction to be read or written. The BM lines are ignored as a bus slave and assume word transfers only. The LANCE drives the BM lines only when it is a bus master. Byte selection occurs as follows:

BM1 BM0

Low Low Whole Word

Low High Byte of DAL 08 - DAL 15

High Low Byte of DAL 00 - DAL 07

High High None

If BCON = 1

PIN 16 = BUSAKO (Output)

PIN 15 = BYTE (Output Three State)

BYTE. An alternate byte selection line. Byte selection occurs when the BYTE and DAL (00) lines are latched during the address portion of the bus transaction. BYTE, $\overline{BM0}$ and $\overline{BM1}$ are ignored when LANCE is a bus slave. There are two modes of ordering bytes depending on bit (02) of CSR3, (known as BSWP). This programmable ordering of upper and lower bytes when using BYTE and DAL (00) as selection signals is required to make the ordering compatible with various 16-bit microprocessors.

BSWP = 0 BSWP = 1

BYTE DAL(00) BYTE DAL(00)

Low Low Low Low Whole Word Low High Low High Illegal Condition High High Hiah Low Upper Byte High Low Lower Byte High High

BUSAKO. The DMA daisy chain output.

HOLD/BUSRQ

(Bus Hold Request)

Input/Output Open Drain. LANCE asserts this signal when it requires access to memory. HOLD is held low for the entire bus transaction. This bit is programmable through bit (00) of CSR3 (known as BCON). In the daisy chain DMA mode (BCON = 1) BUSRQ is asserted only if BUSRQ is inactive prior to assertion. Bit (00) of CSR3 is cleared when RESET is asserted.

CSR3(00) BCON = 0

PIN 17 = HOLD (Output Open Drain)

CSR3(00) BCON = 1

PIN 17 = BUSRQ (Output Open Drain)

BUSRQ will be asserted only if PIN 17 is high prior to assertion.

ALE/AS

(Address Latch Enable)

Output Three State. Used to demultiplex the DAL lines and define the address portion of the bus cycle. This

pin is programmable through bit (01) of CSR3. As ALE, the signal transitions from high to low at the end of the address portion of the bus the address portion of the bus transaction and remains low during the entire data portion of the transaction. As \overline{AS} , the signal transitions from low to high at the end of the address portion of the bus transaction and remains high throughout the entire data portion of the transaction. The LANCE drives the ALE/ \overline{AS} line only as a bus master.

CSR3(01) ACON = 0 PIN 18 = ALE CSR3(01) ACON = 1 PIN 18 = AS

HLDA

(Bus Hold Acknowledge)

Input. A response to HOLD indicating that the LANCE is the Bus Master. HLDA stops its response when HOLD ends its assertion.

cs

(Chip Select)

Input. When asserted, \overline{CS} indicates LANCE is the slave device of the data transfer. \overline{CS} must be valid throughout the data portion of the bus cycle.

ADR

(Register Address Port Select)

Input. When \overline{CS} is asserted, ADR indicates which of the two register ports is selected. ADR must be valid throughout the data portion of the bus cycle.

ADR PORT
Low Register Data Port
High Register Address Port

READY

Input/Output Open Drain. When the LANCE is a bus master, READY is an asynchronous acknowledgement from external memory that will complete the data transfer. As a bus slave, the chip asserts READY when it has put data on the bus, or is about to take data off the bus. READY is a response to DAS. READY negates after DAS negates. Note: If DAS or CS deassert prior to the assertion of READY.

RESET

Input. Bus reset signal. Causes LANCE to cease operation and to enter an idle state.

TLCK

(Transmit Clock)

Input. Normally a free-running 10 MHz clock (crystal-controlled within 0.01% accuracy).

TENA

(Transmit Enable)

Output. Transmit Output Stream Enable. A level asserted with the transmit output bit stream, TX, to enable the external transmit logic.

RCLK

(Receive Clock)

Input. Normally a 10 MHz square wave synchronized to the receive data and present only while receiving an input bit stream.

CLSN

(Collision)

Input. A logical input that indicates that a collision is occurring on the channel.

TX

(Transmit)

Output. Transmit output bit stream.

RENA

(Receive Enable)

Input. A logical input that indicates the presence of data on the channel.

RX

(Receive)

Input. Receive input bit stream.

A16-A23

(High-Order Address Bus)

Output Three State. The additional address bits necessary to extend the DAL lines to produce a 24-bit address. These lines will be driven only as a bus master.

vcc

Power supply pin. +5 VDC ± 5 percent. It is recommended that a power supply filter be used between V_{CC} (Pin 48) and V_{SS} (Pins 1 and 24). This filter should consist of two capacitors in parallel having the values of 10 μ F and .047 μ F respectively.

VSS

Ground, 0 VDC.

FUNCTIONAL CAPABILITIES

The LANCE interfaces to a microprocessor bus characterized by time-multiplexed address and data lines. Typically, data transfers are 16 bits wide but byte transfers occur if the buffer memory address boundaries are odd. The address bus is 24 bits wide.

The Ethernet packet format consists of 64-bit preamble, a 48-bit destination address, a 48-bit source address, a 16-bit type field, and from a 46 to 1500 byte data field terminated with a 32-bit CRC. The packets' variable widths accommodate both short-status command and terminal traffic packets and long data packets to printers and disks (1024-byte disk sectors, for example). Packets are spaced a minimum of 9.6 $\mu \rm sec$ apart to allow one node enough time to receive back-to-back packets.

The LANCE operates in a minimal configuration that requires close coupling between local memory and a processor. The local memory provides packet buffering for the chip and serves as a communication link between chip and processor. During initialization, the control processor loads the starting address of the initialization block plus the operating mode of the chip via two ports that can access four control registers into

LANCE. The host processor talks directly to LANCE only during this initial phase. All further communications are handled via a Direct Memory Access (DMA) machine under microword control contained within LANCE. Figure 3 shows a block diagram of the LANCE and SIA device used to create an Ethernet interface for a computer system.

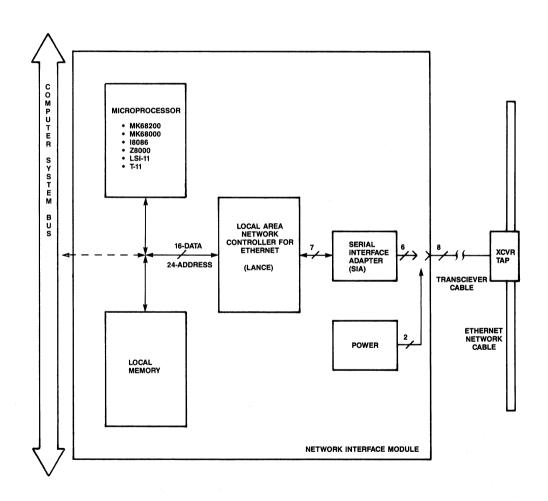


Figure 3. Ethernet Local Area Network System Block Diagram

FUNCTIONAL DESCRIPTION

SERIAL DATA HANDLING

LANCE provides the Ethernet interface as follows. In the transmit mode (since there is only one transmission path, Ethernet is a half duplex system), the LANCE reads data from a transmit buffer by using DMA and appends the preamble, sync pattern (two ones after alternating ones and zeros in the preamble), and calculates and appends the complement of the 32-bit CRC. In the receive mode, the destination address, source address, type, data, and CRC fields are transferred to memory via DMA cycles. The CRC is calculated as data and transmitted CRC is received. At the end of the packet, if this calculated CRC does not agree with a constant, an error bit is set in RDM1 of the receiver descriptor ring. In the receive mode, LANCE accepts packets under four modes of operation. The first mode is a full comparison of the 48-bit destination address in the packet with the node address that was programmed into the LANCE during an initialization cycle. There are two types of logical addresses. One is a group type mask where the 48-bit address in the packet is put through a hash filter in order to map the 48-bit physical addresses into 1 of 64 logical groups. This mode can be useful if simultaneously sending packets to all of one type of a device on the network, (i.e., send a packet to all file servers or all printer servers). The second logical address is a multicast address where all nodes on the network receive the packet. The last receive mode of operation is the so called "promiscuous mode" in which a node will accept all packets on the cable regardless of their destination address.

COLLISION DETECTION AND IMPLEMENTATION

The Ethernet CSMA/CD network access algorithm is implemented completely within LANCE. In addition to listening for a clear network cable before transmitting, Ethernet handles collisions in a predetermined way. Should two transmitters attempt to seize the network cable at the same time, they will collide, and the data on the network cable will be garbled. LANCE is constantly monitoring the Collision (CLSN) pin. This signal is generated by the transceiver when the signal level on the network cable indicates the presence of signals from two or more transmitters. If LANCE is transmitting when CLSN is asserted, it will continue to transmit the preamble (collisions normally occur while the preamble is being transmitted), then will "jam" the network for 32 bit times (3.2 microseconds). This jamming ensures that all nodes have enough time to detect the collision. The transmitting nodes then delay a random amount of time according to the "truncated binary backoff" algorithm defined in the Ethernet specification to minimize the probability of the colliding nodes having multiple collisions with each other. After 16 abortive attempts to transmit a packet, LANCE will report a RTRY error due to excessive collisions and step over the transmitter buffer. During reception, the detection of a collision causes that reception to be aborted. Depending on when the collision occurred, LANCE will treat this packet as an error packet if the packet has an address mismatch, as a runt packet (a packet that has less than 64 bytes), or as a legal length packet with a CRC error.

Fatal error reporting is provided by LANCE through a microprocessor interrupt and error flags in CSRO. These Fatal Error conditions are the following:

- Collision Error The failure of the transceiver to send a collision message or, heartbeat at the conclusion of a normal transmission.
- 2.) Transmitter on longer than 1518 bytes.
- Missed Packet LANCE failed to receiver transmitted packet.
- 4.) Memory Error Failure of a memory transaction to complete within 25.6 μ S.

Additional errors are reported through bits in the descriptor rings (on a buffer by buffer basis). Receive error conditions include framing, CRC and buffer errors, and overflow. Transmit descriptor rings have error bits indicating buffer, underflow, late collision, and loss of carrier. Additionally, transmit descriptor rings have a bit indicating that the transmitter has unsuccessfully tried to transmit over a busy communication link.

Transmit descriptor rings also have ten bits reserved for a Time Domain Reflectometry counter (TDR). On the occurrance of a collision, the value in the TDR will give the number of system clocks until the collision, which can be used to determine the distance to the fault.

BUFFER MANAGEMENT

A key feature of the LANCE and its DMA channel is the flexibility and speed of communication between the LANCE and the host microprocessor through common memory locations. The basic organization of the buffer management is a circular queue of tasks in memory called descriptor rings, as shown in Figure 4. These rings control both transmit and receive operations. Up to 128 tasks may be gueued on a descriptor ring for execution by the LANCE. Each entry in a descriptor ring holds a pointer to a data memory buffer and an entry for the data buffer length. Data buffers can be chained or cascaded to handle a long packet in multiple data buffer areas. The LANCE searches the descriptor rings to determine the next empty buffer. This enables it to chain buffers together or to handle back-to-back packets. As each buffer is filled, an "own" bit is reset, signaling the host processor to empty this buffer.

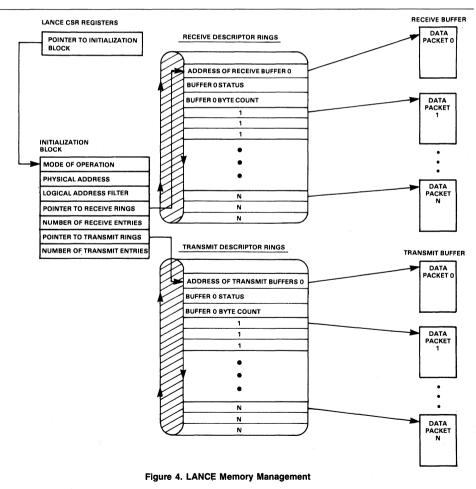
MICROPROCESSOR INTERFACE

The parallel interface of LANCE has been designed to be "friendly", or easy to interface, to many popular 16-bit microprocessors. These microprocessors include the MK68000, Z8000, 8086, LSI-11, T-11, and MK68200 (the MK68200 is a 16-bit single chip microcomputer produced by Mostek, the architecture of which is modeled after the MK68000). LANCE has a wide 24-bit linear address space when in the Buster Master Mode, allowing it to DMA the entire address space of the above microprocessors. LANCE uses no segmentation or paging methods. As such, LANCE addressing is closest to MK68000 addressing, but is compatible with the other microprocessors. When LANCE is a bus master. a programmable mode of operation allows byte addressing, either by employing a Byte/Word control signal (much like that used on the 8086 or the Z8000)

or by using an Upper Data Strobe/Lower Data Strobe much like that used on the MK68000, LSI-11 and MK68200 microprocessors. A programmable polarity on the Address Strobe signal eliminates the need for external logic. LANCE interfaces with multiplexed and demultiplexed data busses and features control signals for address/data bus transceivers.

After the initialization routine, packet reception or transmission, transmitter timeout error, a missed packet, or memory error, LANCE generates an interrupt to the host microprocessor.

The cause of the interrupt is ascertained by reading CSR0. Bit (06) of CSR0, INEA, enables or disables interrupts to the microprocessor. In a polling mode, BIT (07) of CSR0 is sampled to determine when an interrupt causing condition occurred.



LANCE INTERFACE DESCRIPTION

ALE, DAS and READY time all data transfers from the LANCE in the Bus Master mode. The automatic adjustment of the LANCE cycle by the READY signal allows synchronization with variable cycle time memory due either to memory refresh or to dual port access. Bus cycles are a minimum of 600 ns long and can be increased in 100 ns increments.

READ SEQUENCE

At the beginning of a read cycle, valid addresses are placed on DAL00-DAL15 and A16-A21. The BYTE Mask signals (BMO and BM1) become valid at the beginning of this cycle as does READ, indicating the type of cycle. The trailing edge of ALE or AS strobes the addresses A0-A15 into the external latches. Approximately 100 ns later, DAL00-DAL15 go into a three state mode. There is a 50 ns delay to allow for transceiver turnaround, then DAS falls low to signal the beginning of the data portion of the cycle. At this point in the cycle, the LANCE stalls waiting for the memory device to assert READY. Upon assertion of READY, DAS makes a transition from a zero to a one, latching memory data. (DAS is low for a minimum of 200 ns).

The bus transceiver controls, \overline{DALI} and \overline{DALO} , control

the bus transceivers. DALI signals to strobe data toward the LANCE and DALO signals to strobe data or addresses away from the LANCE. During a read cycle, DALO goes inactive before DALI goes active to avoid "spiking" of bus transceivers.

WRITE SEQUENCE

The write cycle begins exactly like a read cycle with the READ line remaining inactive. After ALE or $\overline{\text{AS}}$ pulse, the DAL00-DAL15 change from addresses to data. $\overline{\text{DAS}}$ goes active when the DAL00-DAL15 are stable. This data remains valid on the bus until the memory device asserts $\overline{\text{READY}}$. At this point, $\overline{\text{DAS}}$ goes inactive, latching data into the memory device. Data is held for 75 ns after the negation of $\overline{\text{DAS}}$.

LANCE INTERFACE DESCRIPTION — BUS SLAVE MODE

The LANCE enters the Bus Slave Mode whenever $\overline{\text{CS}}$ becomes active. This mode must be entered whenever writing or reading the four status control registers (CSR0, CSR1, CSR2, and CSR3) and the register address pointer (RAP). RAP and CSR0 may be read or written to at any time, but the LANCE must be stopped (CSR0 bit 02) when CSR1, CSR2, or CSR3 is to be written to or read.

MK68590 ELECTRICAL SPECIFICATION

ABSOLUTE MAXIMUM RATINGS

Temperature Under Bias25°C to +125°C
Storage Temperature65°C to +150°C
Voltage on Any Pin with Respect to Ground
Power Dissipation

Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other condition above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

DC CHARACTERISTICS

 $T_A=0$ °C to 70 °C, $V_{CC}=+5$ V ± 5 percent unless otherwise specified.

SYMBOL	PARAMETER	MIN	MAX	UNITS
V _{IL}		-0.5	+0.8	٧
V _{IH}		+2.0	V _{CC} +0.5	٧
V _{OL}	$@ I_{OL} = 3.2 \text{ mA}$		+0.5	٧
V _{OH}	$@ I_{OH} = -0.4 \text{ mA}$	+2.4		٧
I _{IL}	$@V_{in} = 0.4 \text{ to } V_{CC}$		±10	μΑ

CAPACITANCE

F=1 MHz

SYMBOL	PARAMETER	MIN	MAX	UNITS
C _{IN}			10	pf
C _{OUT}			10	pf
C _{IO}			20	pf

AC TIMING SPECIFICATIONS

 T_A = 0 °C to 70 °C, V_{CC} = +5 V ± 5 percent, unless otherwise specified.

NO.	SIGNAL	SYMBOL	PARAMETER	TEST CONDITIONS	MIN (ns)	TYP (ns)	MAX (ns)
1	TCLK	T _{TCT}	TCLK period		99		101
2	TCLK	T _{TCL}	TCLK low time		45		55
3	TCLK	T _{TCH}	TCLK high time		45		55
4	TCLK	T _{TCR}	Rise time of TCLK		0		8
5	TCLK	T _{TCF}	Fall time of TCLK		0		8
6	TENA	T _{TEP}	TENA propagation delay after the rising edge of TCLK	CL = 50 pf			75
7	TENA	T _{TEH}	TENA hold time after the rising edge of TCLK	CL = 50 pf	5		

AC TIMING SPECIFICATIONS (Continued) $T_A = 0\,^{\circ}\!C \ to \ 70\,^{\circ}\!C, \ V_{CC} = +5 \ V \ \pm 5 \ percent, \ unless \ otherwise \ specified.$

NO.	SIGNAL	SYMBOL	PARAMETER	TEST CONDITIONS	MIN (ns)	TYP (ns)	MAX (ns)
8	TX	T _{TDP}	TX data propagation delay after the rising edge of TCLK	CL=50 pf			75
9	TX	T _{TDH}	TX data hold time after the rising edge of TCLK	CL=50 pf	5		
10	RCLK	T _{RCT}	RCLK period		85		118
11	RCLK	T _{RCH}	RCLK high time		38		
12	RCLK	T _{RCL}	RCLK low time		38		
13	RCLK	T _{RCR}	Rise time of RCLK		0		8
14	RCLK	T _{RCF}	Fall time of RCLK		0		8
15	RX	T _{RDR}	RX data rise time		0		8
16	RX	T _{RDF}	RX data fall time		0		8
17	RX	T _{RDH}	RX data hold time (RCLK to RX data change)		5		
18	RX	T _{RDS} (See Note)	RX data setup time (RX data stable to the rising edge of RCLK)		See Note		
19	RENA	T _{DPL}	RENA low time		120		
20	RENA	T _{RENH}	RENA hold time after rising edge of RCLR		40		
21	CLSN	T _{CPH}	CLSN high time		80		
22	A/DAL	T _{DOFF}	Bus master driver disable after rising edge of HOLD		0		50
23	A/DAL	T _{DON}	Bus master driver enable after falling edge of HLDA		0		150
24	HLDA	T _{HHA}	Delay to falling edge of HLDA from falling edge of HOLD (bus master)		0		
25	RESET	T _{RW}	RESETpulse width low		200		
26	A/DAL	T _{CYCLE}	Read/write, address/data cycle time		600		
27	Α	T _{XAS}	Address setup time to the falling edge of ALE		75		
28	Α	T _{XAH}	Address hold time after the rising edge of $\overline{\text{DAS}}$		35		
29	DAL	T _{AS}	Address setup time to the falling edge of ALE		75		
30	DAL	T _{AH}	Address hold time after the falling edge of ALE		35		
31	DAL	T _{RDAS}	Data setup time to the rising edge of DAS (bus master read)		50		

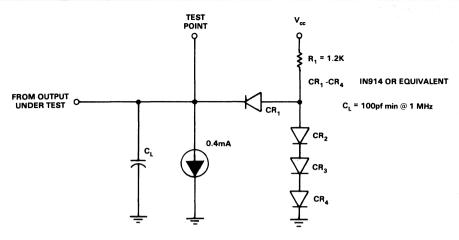
NOTE: $T_{RDS \ (min)} = T_{RCT} - 25 \text{ ns. Therefore, } T_{RCT} = 100 \text{ ns when } T_{RDS \ (min)} = 75 \text{ ns.}$

AC TIMING SPECIFICATIONS (Continued)

 $T_A = 0$ °C to 70°C, $V_{CC} = +5$ V ± 5 percent, unless otherwise specified.

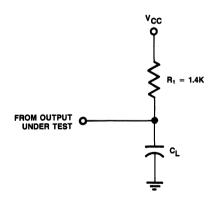
NO.	SIGNAL	SYMBOL	PARAMETER	TEST CONDITIONS	MIN (ns)	TYP (ns)	MAX (ns)
32	DAL	T _{RDAH}	Data hold time after the rising edge of DAS (bus master read)		0		
33	DAL	T _{DDAS}	Data setup time to the falling edge of DAS (bus master write)		0		
34	DAL	T _{WDS}	Data setup time to the rising edge of DAS (bus master write)		200		
35	DAL	T _{WDH}	Data hold time after the rising edge of DAS (bus master write)		35		
36	DAL	T _{SDO1}	Data driver delay after the falling edge of DAS (bus slave read)	(CSR 0,3, RAP)		400	
37	DAL	T _{SDO2}	Data driver delay after the falling edge of DAS (bus slave read)	(CSR 1,2)		1200	
38	DAL	T _{SRDH}	Data hold time after the rising edge of DAS (bus slave read)		0		35
39	DAL	T _{SWDH}	Data hold time after the rising edge of DAS (bus slave write)		0		
40	DAL	T _{SWDS}	Data setup time to the falling edge of DAS (bus slave write)		0		
41	ALE	T _{ALEW}	ALE width high		120		150
42	ALE	T _{DALE}	Delay from rising edge of DAS to the rising edge of ALE		70		
43	DAS	T _{DSW}	DAS width low		200		
44	DAS	T _{ADAS}	Delay from the falling edge of ALE to the falling edge of DAS		80		130
45	DAS	T _{RIDF}	Delay from the rising edge of DALO to the falling edge of DAS (bus master read)		15		
46	DAS	T _{RDYS}	Delay from the falling edge of READY to the rising edge of DAS	Taryd= 300 ns	75		250
47	DALI	T _{ROIF}	Delay from the rising edge of DALO to the falling edge of DALI (bus master read)		15		
48	DALI	T _{RIS}	DALI setup time to the rising edge of DAS (bus master read)		135		
49	DALI	T _{RIH}	DALI hold time after the rising edge of DAS (bus master read)		0	-	
50	DALI	T _{RIOF}	Delay from the rising edge of DALI to the falling edge of DALO (bus master read)		55		
51	DALO	Tos	DALO setup time to the falling edge of ALE (bus master read)		110		
52	DALO	T _{ROH}	DALO hold time after the falling edge of ALE (bus master read)		35		
53	DALO	T _{WDSI}	Delay from the rising edge of DAS to the rising edge of DALO (bus master write)		35		

NO.	SIGNAL	SYMBOL	PARAMETER	TEST CONDITIONS	MIN (ns)	TYP (ns)	MAX (ns)
54	CS	T _{CSH}	CS hold time after the rising edge of DAS (Bus slave)		0		
55	CS	T _{CSS}	CS setup time to the falling edge of DAS (Bus slave)		0		
56	ADR	T _{SAH}	ADR hold time after the rising edge of DAS (Bus-slave)		0		
57	ADR	T _{SAS}	ADR setup time to the falling edge of DAS (Bus slave)		0		
58	READY	T _{ARYD}	Delay from the falling edge of ALE to the falling edge of READY to insure a minimum bus cycle time (600 ns)				80
59	READY	T _{SRDS}	Data setup time to the falling edge of READY (Bus slave read)		75		
60	READY	T _{RDYH}	READY hold time after the rising edge of DAS (Bus master)		0		
61	READY	T _{SRO1}	READY driver turn on after the falling edge of DAS (Bus slave)	(CSR 0,3, RAP)	-	600	
62	READY	T _{SRO2}	READY driver turn on after the falling edge of DAS (Bus slave)	(CSR 1,2)		1400	
63	READY	T _{SRYH}	READY hold time after the rising edge of DAS (Bus slave)		0		35
64	READ	T _{SRH}	READ hold time after the rising edge of \overline{DAS} (Bus slave)		0		
65	READ	T _{SRS}	READ setup time to the falling edge of DAS (Bus slave)		0		



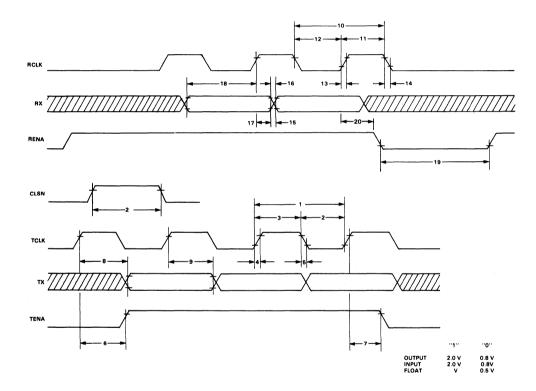
NOTE: This load is used on DAL00 through DAL15, READ, DAL1, DALO, DAS, BM0, BM1, ALE/AS, A16 through A23, TENA, and TX.

Figure 5. Output Load Diagram



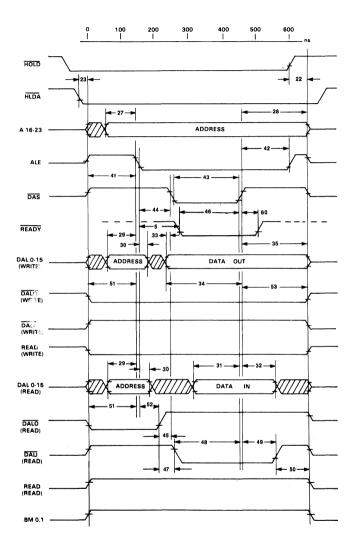
NOTE: This load is used on open drain outputs $\overline{\text{INTR}}$, HOLD/BUSRQ, and $\overline{\text{READY}}$.

Figure 6. Open Drain Output Load Diagram.



NOTE: Timing measurements are made at the following voltages unless otherwise specified.

Figure 7. Serial Link Timing Diagram - SIA Interface Signals



NOTE: The Bus Master cycle time will increase from a minimum of 600 ns in 100 ns steps until the slave device returns READY.

Figure 8. LANCE Bus Master Timing Diagram

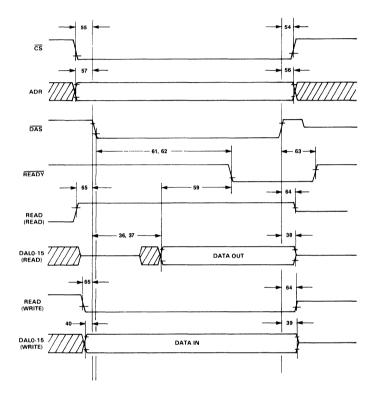


Figure 9. LANCE Bus Slave Timing Diagram